

AMAT 584 Lecture 37 4/27/20

Today: The persistence algorithm, continued

Review: To compute the persistent homology of a filtration F with largest simplicial complex F_{\max} :

1) For each $j \geq 0$, order simplices of F_{\max} with respect to their birth time in the filtration.

2) Using these orderings, construct a "boundary matrix" D containing all the matrices $[\delta_j]$ as blocks just above the diagonal.

Columns and rows of D correspond to simplices of F_{\max} .

3) Use left-to-right column additions to transform D into a reduced matrix R .

4) Read each barcode off of R , using the following formula:

$\text{Barc}(H_i(F)) =$

$\{[\text{birth}(\sigma), \text{birth}(\tau)] \mid \text{pivot of column } \tau \text{ in } R \text{ is } \sigma \text{ and } \dim(\sigma) = i\}$

$\{ [birth(\sigma), \infty) \mid col \sigma = 0, \sigma \text{ is not the pivot of any column in } R, dim(\sigma) = i \}$.

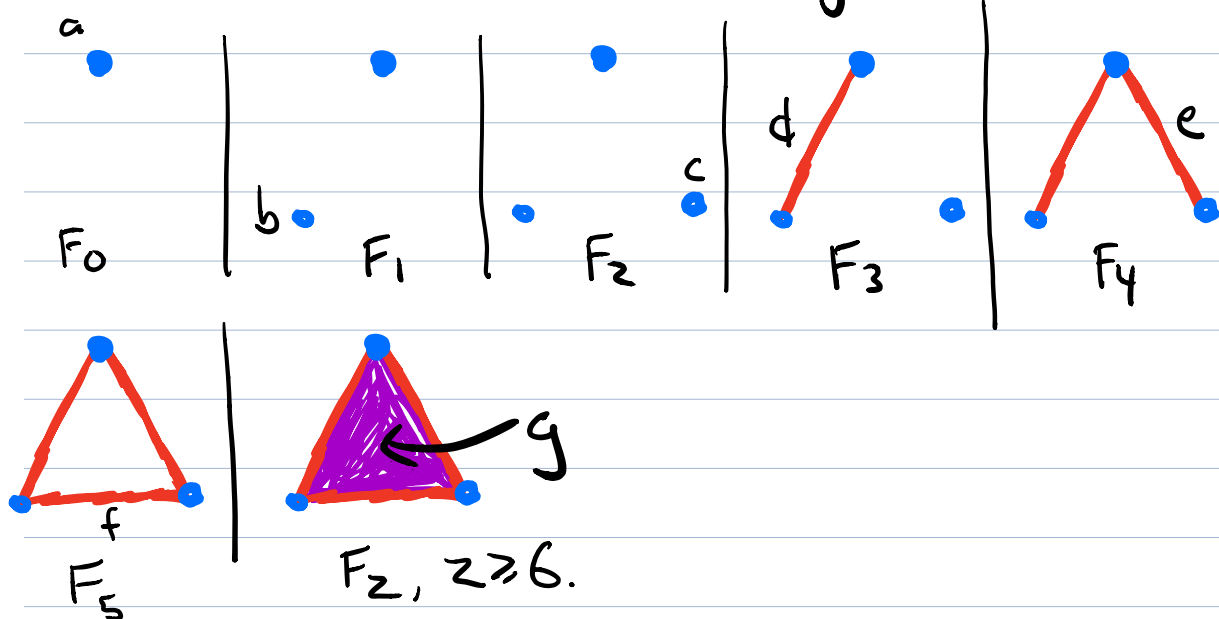
Note: Any interval the form $[z, z)$ is ignored.

In words: Every non-zero column in R gives a finite-length interval in one of the barcodes, by pairing the column with its pivot.

A zero column which doesn't appear as the pivot of any other column gives an infinite interval.

Example of the last step

Last lecture we considered the following filtration F :



We saw that for this filtration,

$$D = \begin{array}{c} a \\ b \\ c \\ d \\ e \\ f \\ g \end{array} \begin{array}{ccccccc} a & b & c & d & e & f & g \\ \hline & & & 1 & 1 & 0 & \\ \hline & 0 & & 1 & 0 & 1 & 0 \\ \hline & & & 0 & 1 & 1 & \\ \hline & & & & & & 1 \\ \hline & 0 & & & 0 & & 1 \\ \hline & & & & & & 1 \\ \hline & 0 & & & 0 & & 0 \end{array}$$

Applying the standard reduction (step 3) gives

$$R = \begin{array}{c} F_{\max}^0 \\ \\ \\ F_{\max}^1 \\ \\ F_{\max}^2 \end{array} \begin{array}{c} a \\ b \\ c \\ d \\ e \\ f \\ g \end{array} \begin{array}{ccccccc} a & b & c & d & e & f & g \\ \hline & & & 1 & 1 & 0 & \\ \hline & 0 & & 1 & 0 & 0 & 0 \\ \hline & & & 0 & 1 & 0 & \\ \hline & & & & & & 1 \\ \hline & 0 & & & 0 & & 1 \\ \hline & & & & & & 1 \\ \hline & 0 & & & 0 & & 0 \end{array}$$

(The only column of D that was changed was column f).

We get a finite interval in one of the barcodes for each non-zero column:

$$\begin{array}{ccc} \begin{array}{c} F_{\max}^0 \\ \cup \\ [\text{birth}(b), \text{birth}(d)) \\ \cup \\ [1, 3) \end{array} & , & \begin{array}{c} F_{\max}^0 \\ \cup \\ [\text{birth}(c), \text{birth}(e)) \\ \cup \\ [2, 4) \end{array} & , & \begin{array}{c} F_{\max}^1 \\ \cup \\ [\text{birth}(f), \text{birth}(g)) \\ \cup \\ [5, 6) \end{array} \end{array}$$

The zero columns are a, b, c, f . Of these, all but a is a pivot of another column.

So the only infinite interval is $\begin{array}{c} F_{\max}^0 \\ \cup \\ [\text{birth}(a), \infty) \\ \cup \\ [0, \infty) \end{array}$

To figure out which interval goes in which barcode, we look at the dimension of the simplex ending the left endpoint.

$$\text{Barc}(H_0(F)) = \{ [1, 3), [2, 4), [0, \infty) \}$$

$$\text{Barc}(H_1(F)) = \{ [5, 6) \}.$$

Note: If we only want to compute persistent homology up to degree k , i.e. $\text{Barc}(H_i(F))$ for $i \in \{0, 1, \dots, k\}$, then we don't need to consider all of F . It's sufficient to consider the subfiltration consisting of only simplices of dimension $\leq k+1$.

This often allows us to keep the size of D much smaller

An example coming from data

Let $X = \{ \overset{a}{(0,0)}, \overset{b}{(2,0)}, \overset{c}{(0,2)}, \overset{d}{(2,2)} \} \subset \mathbb{R}^2$, given the Euclidean metric.

Let $B_i = \text{Barc}(H_i(\text{Rips}(X)))$

We will compute each B_i .

Note that $\text{Rips}(X)$ is a $[0, \infty)$ -index filtration. We will show by example that the persistence algorithm works in essentially the same way for such a filtration, as in the \mathbb{N} -indexed case.

Let $F = \text{Rips}(X)$. F_{\max} is the 3-simplex with vertex set X .

Let's give names to all the simplices of F_{\max} , and record their birth times:

to keep notation simple, we write $\text{birth}(\sigma)$ as $\beta(\sigma)$.

$$a = [a], \quad b = [b], \quad c = [c], \quad d = [d]$$

$$\beta(a) = \beta(b) = \beta(c) = \beta(d) = 0$$

$$e=[a,b], f=[b,c], g=[c,d], h=[d,a], i=[a,c], j=[b,d]$$

$$\beta(e)=\beta(f)=\beta(g)=\beta(h)=2.$$

$$\beta(i)=\beta(j)=2\sqrt{2}.$$

$$k=[a,b,c], l=[a,b,d], m=[a,c,d], n=[b,c,d]$$

$$\beta(k)=\beta(l)=\beta(m)=\beta(n)=2\sqrt{2}$$

$$o=[a,b,c,d]$$

$$\beta(o)=2\sqrt{2}.$$

Note: The alphabetical order on each X^i is in order of increasing birth

to be continued: